A detailed proof of the Monotone Convergence Theorem My thoughts on writing: Think through writing. Learn through writing. Write, for the andrence of one: May 19,2020 fourself.

The formulation 23 Learn Horough writing of fle Competences Afrian below is White for yourself? Zorich The Completeners Axions of TR \* If A = R, B = R, A + Ø, B + Ø and FAGA HEB we have a < b, then FCER s.t. adeceb Hacatber A <u>Brumber line</u> IR

The Monotone Convergence Theorem If a sequence is monotonic and bounded, then it converges. Proof. Case 1. A signence is non-decreasing and bounded above Let A: NOR be non decreasing and bold above That is  $\forall n \in \mathcal{N}$   $\Lambda_n \leq \Lambda_{n+1} (\Lambda_1 \leq \Lambda_2 \leq \dots \leq \Lambda_n \leq \Lambda_{n+1} \leq \dots$ and J MER such that the IN Sn EM What is RED? JLERS. H. HE>D JN(E)ERS. T. FREAN we have n=N(e)=) Sn-L/<E

S1 S2 S2 S4 Dn Dn+1 M1 M M2 TR This smells like CA? (Completeness Axion) So, set A = { In : n e M } (range of J:NR) A = { In : n e M } (range of J:NR) B = JbER: bis an upper bound for Je Clearly A7Ø, B7Ø since MEB. Clearly Hack Hbe B we have  $a \leq b$ I a  $\in A$ , then  $a = J_n \leq b \in B$ since  $b \leq a$  upper n  $\in X$ ?

Thus All the hypothesis of CA are satisfied Therefore JCER such that YNENTBEB Snéréb G1 any dement of A A By G1 C is an upper bound for the sequence s. my suppor bound Thus  $c \in B$  (all upper bounds for  $\Lambda$ ). (c is called the least upper bound for S) In fact we have that  $c = \min B$ C is the minimum of B

Redo our mbor live with A and B  $\mathbb{R}$  $\frac{\Delta_4}{\Delta_1} \frac{\Delta_{n+1}}{\Delta_2} \frac{b_1}{\delta_2} \frac{b_2}{\delta_1} \frac{b_2}{\delta_1} \frac{b_1}{\delta_1} \frac{b_2}{\delta_1} \frac{b_1}{\delta_1} \frac{b_2}{\delta_2} \frac{b_1}{\delta_1} \frac{b_1}{\delta_2} \frac{b_2}{\delta_1} \frac{b_1}{\delta_1} \frac{b_2}{\delta_2} \frac{b_1}{\delta_1} \frac{b_1}{\delta_1} \frac{b_2}{\delta_1} \frac{b_1}{\delta_1} \frac{b_1}{\delta_1} \frac{b_2}{\delta_1} \frac{b_1}{\delta_1} \frac{b_1}{\delta_1} \frac{b_2}{\delta_1} \frac{b_1}{\delta_1} \frac{b_1}{\delta_1}$ A is special it is the last upper bound for this NOT an upper bound. Now we are ready to address REDT Set L= construction. Let 2>0 be arbitrary. Then c-E < c, but  $c \leq b$  for all b upper bounds for s. Therefore c-Eris NOT an upper bound forso

C-E is NOT an upper bound for S. How do you say this in Mathish language? This is trow: ZN(E) EN such that C-E<AN(E) G2 Now if I take n > N(e)since  $M_n$  in non-decreasing, I have  $G3 M(e) \leq Mn > N(e)$ 

From G2 & G3 we conclude that  $\forall n \in \mathbb{N}$   $n \geq \mathbb{N}(\varepsilon) \Rightarrow \mathcal{L} - \varepsilon \leq \mathcal{I}_n.$ That is: Thek n>N(E) = C-Sn < EG4 But we know from G1 that the N C>Sn. Therefore the N [Sn-C] = C-Sn. Thus G4 can be rewritten as  $\forall n \in \mathbb{N} \quad n \geq \mathbb{N}(\varepsilon) \Rightarrow | \Lambda_n - c | < \varepsilon$ 

Since E>O was arbitrary, we have proved that HE>O JN(E)EN Such that  $\forall n \in \mathbb{N}$   $n \geq \mathbb{N}(\varepsilon) \Rightarrow | \mathbb{A}_{-\mathcal{C}} | < \varepsilon$ . This proves that L = c is the limit of  $S: N \rightarrow R$ . We proved  $\lim_{n \to +\infty} \Lambda_n = \mathcal{L}_{\bullet}$